

# Beyond the Embarrassingly Parallel

New Languages, Compilers, and Runtimes for Big-Data Processing

Madan Musuvathi

Microsoft Research

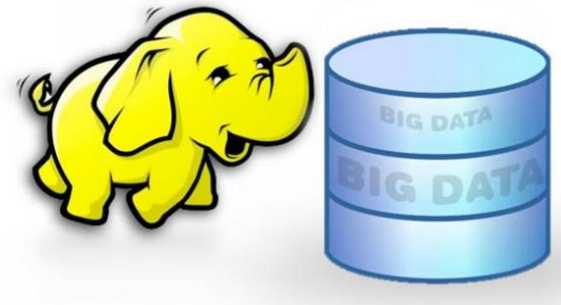
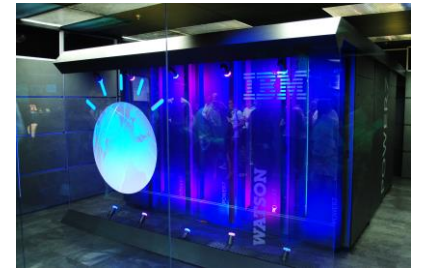
Joint work with

Mike Barnett (MSR), Saeed Maleki (MSR), Todd Mytkowicz (MSR)

Yufei Ding (N.C.State), Daniel Lupei (EPFL), Charith Mendis (MIT),

Mathias Peters (Humboldt Univ.), Veselin Raychev (EPFL)

# parallelism



parallelism  
=  
independent computation

can we parallelize  
dependent computation?

# “Inherently sequential” code is common



log processing

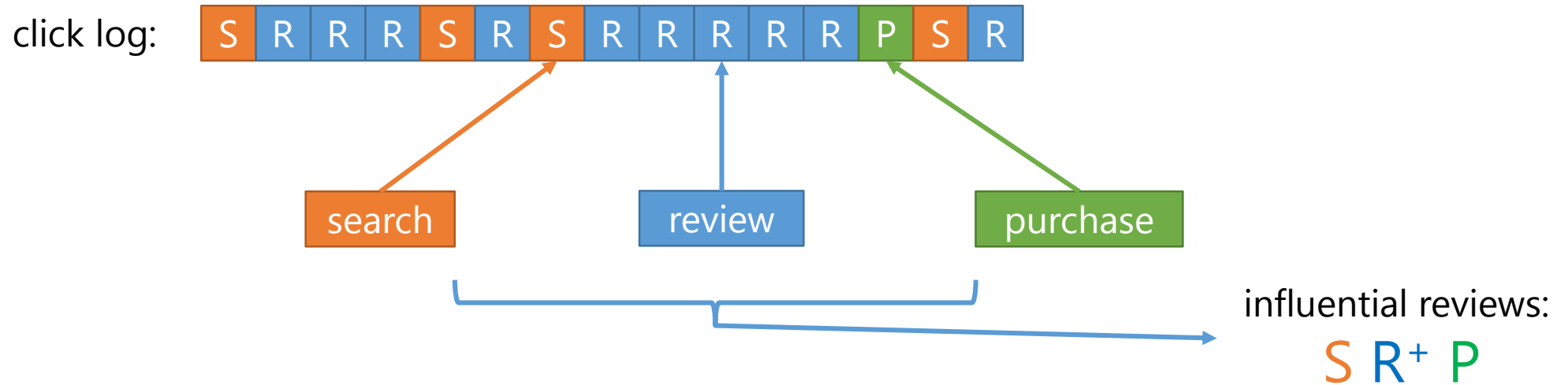
event-series pattern matching

machine learning algorithms

dynamic programming

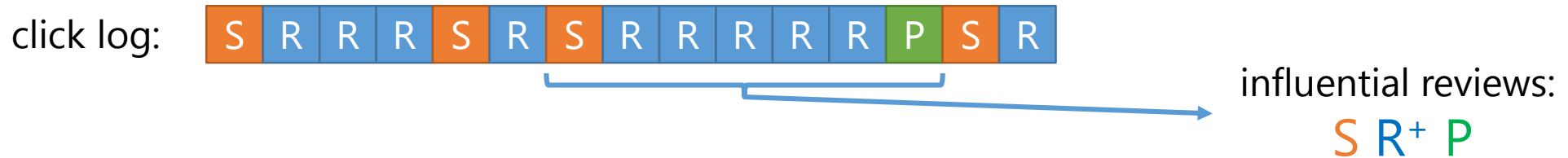
...

# Running example: processing click logs



problem: count influential reviews in the log

# Running example: processing click logs



```
bool search_done = false; int num_reviews = 0; int sum = 0;
```

```
for each record in input
  switch record.type:
    case SEARCH:    if (!search_done) { num_reviews = 0;
                                     search_done = true; }
    case REVIEW:    num_reviews++;
    case PURCHASE:  if (search_done) { search_done = false;
                                     sum += num_reviews; }
```

Loop carried state

# Extracting parallelism from dependent computations



```
for each record in input
  switch record.type:
    case SEARCH:  if (!search_done) { num_reviews = 0;
                                     search_done = true; }
    case REVIEW:  num_reviews++;
    case PURCHASE: if (search_done) { search_done = false;
                                     sum += num_reviews; }
```

(true, 1, 2)

```
for each record in input
  switch record.type:
    case SEARCH:  if (!search_done) { num_reviews = 0;
                                     search_done = true; }
    case REVIEW:  num_reviews++;
    case PURCHASE: if (search_done) { search_done = false;
                                     sum += num_reviews; }
```

(false, 0, 0)

(false, 1, 8)

// loop-carried state:  
// (search\_done, num\_reviews, sum)



# Extracting parallelism from dependent computations



```
for each record in input
  switch record.type:
    case SEARCH:  if (!search_done) { num_reviews = 0;
                                     search_done = true; }
    case REVIEW:  num_reviews++;
    case PURCHASE: if (search_done) { search_done = false;
                                     sum += num_reviews; }
```

(false, 0, 0)

// loop-carried state:  
// (search\_done, num\_reviews, sum)

(true, 1, 2)

(sd, nr, s)



```
for each record in input
  switch record.type:
    case SEARCH:  if (!search_done) { num_reviews = 0;
                                     search_done = true; }
    case REVIEW:  num_reviews++;
    case PURCHASE: if (search_done) { search_done = false;
                                     sum += num_reviews; }
```

summary = F(sd, nr, s)

$F(sd, nr, s) = (\text{false}, nr+6, sd ? s+nr+5 : s)$

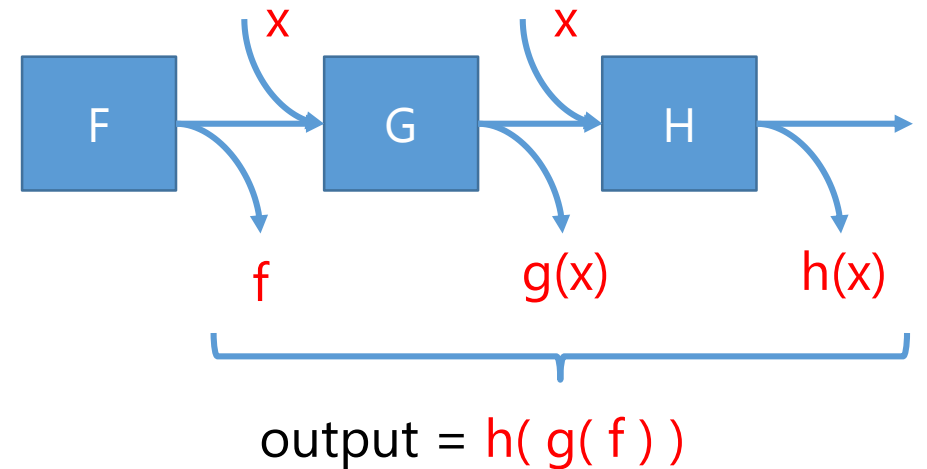
output = F(true, 1, 2)  
= (false, 1, 8)

# Recipe for breaking dependences

1. replace dependences with symbolic unknowns
2. compute symbolic summaries in parallel
3. combine symbolic summaries

success depends on

1. fast symbolic execution
2. generation of concise summaries



research challenges :

1. identifying "compressible" computation
2. using domain-specific structure
3. automating the parallelization

# Successful applications of this methodology

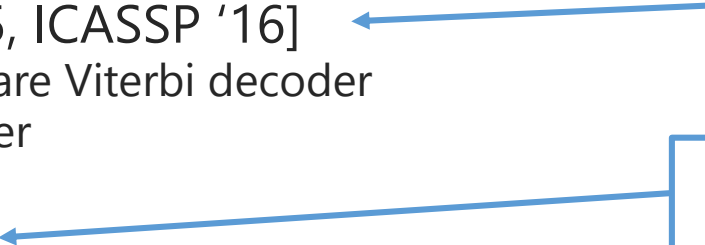
## finite-state machines [ASPLOS '14]

- regular expression matching, Huffman decoding, ...
- 3x faster on a single core, linear speedup on multiple cores

## dynamic programming [PPoPP '14, TOPC '15, ICASSP '16]

- linear speedup beyond the previous-best software Viterbi decoder
- 7x speedup over state-of-the-art speech decoder

part 2 of the talk



## large-scale data processing [SOSP '15]

- automatically parallelizable language for temporal analysis

part 1 of the talk

## relational databases

- optimize sessionization & windowed aggregates
- 10x improvement over SQL server

## machine learning

- parallel stochastic gradient descent

# Auto-Parallelization Across Dependences

Large-scale data processing

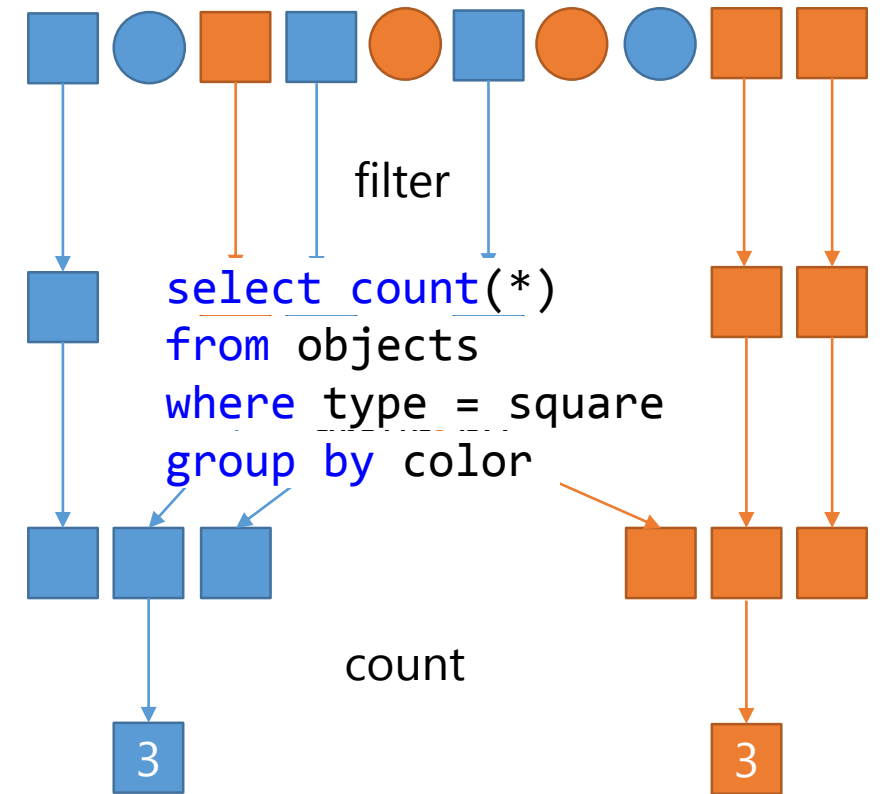
# Relational abstractions for data processing

map, reduce, join, filter, group-by

expressive, simple, and declarative

automatically parallelizable

decades of work on optimizations



# Forces pushing beyond relational abstractions

queries today = relational skeleton + non-relational logic

embarrassingly parallel  
optimized

not parallel  
not optimized

temporal, iterative, stateful

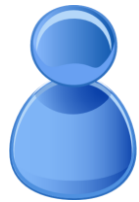
- log analysis
- sessionization
- machine learning

# Map-Reduce example

weblog



users can:



S

R

P

search

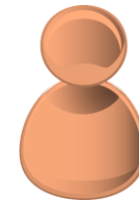
review

purchase

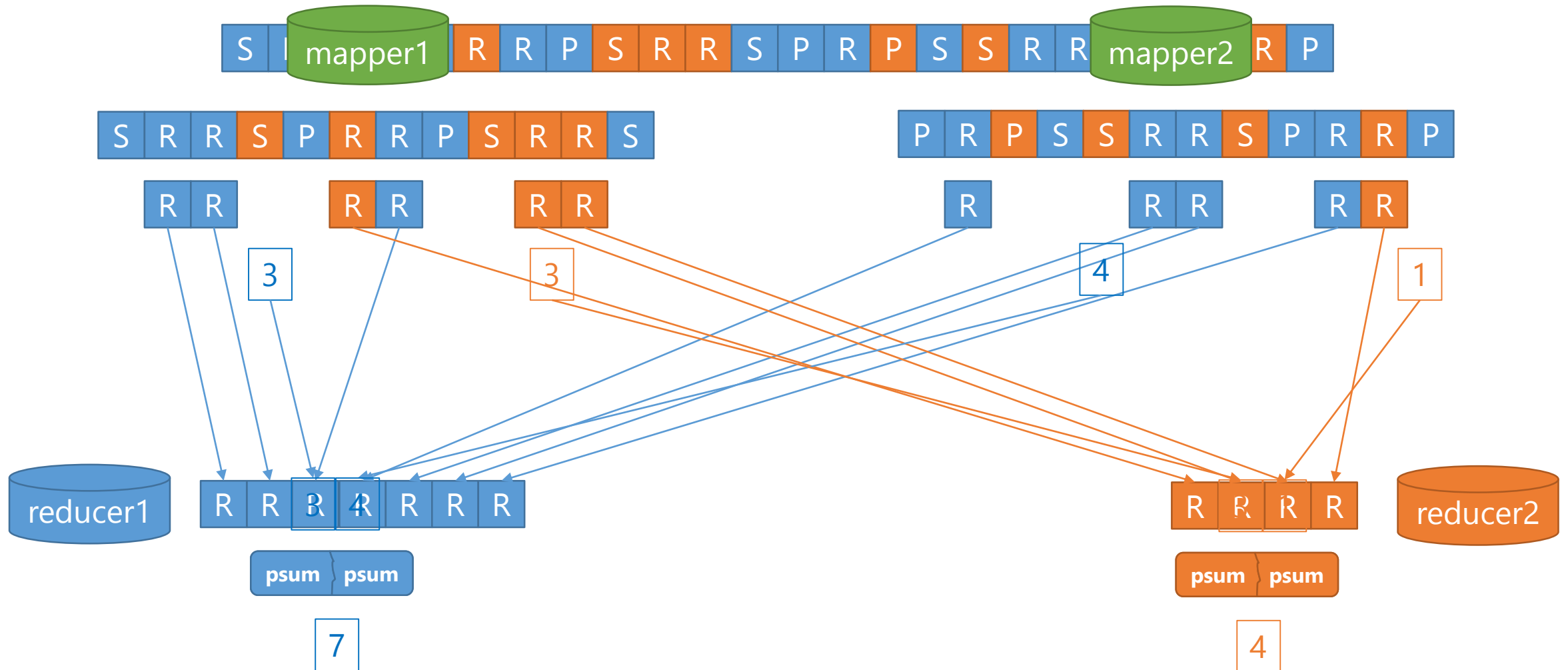
S

R

P

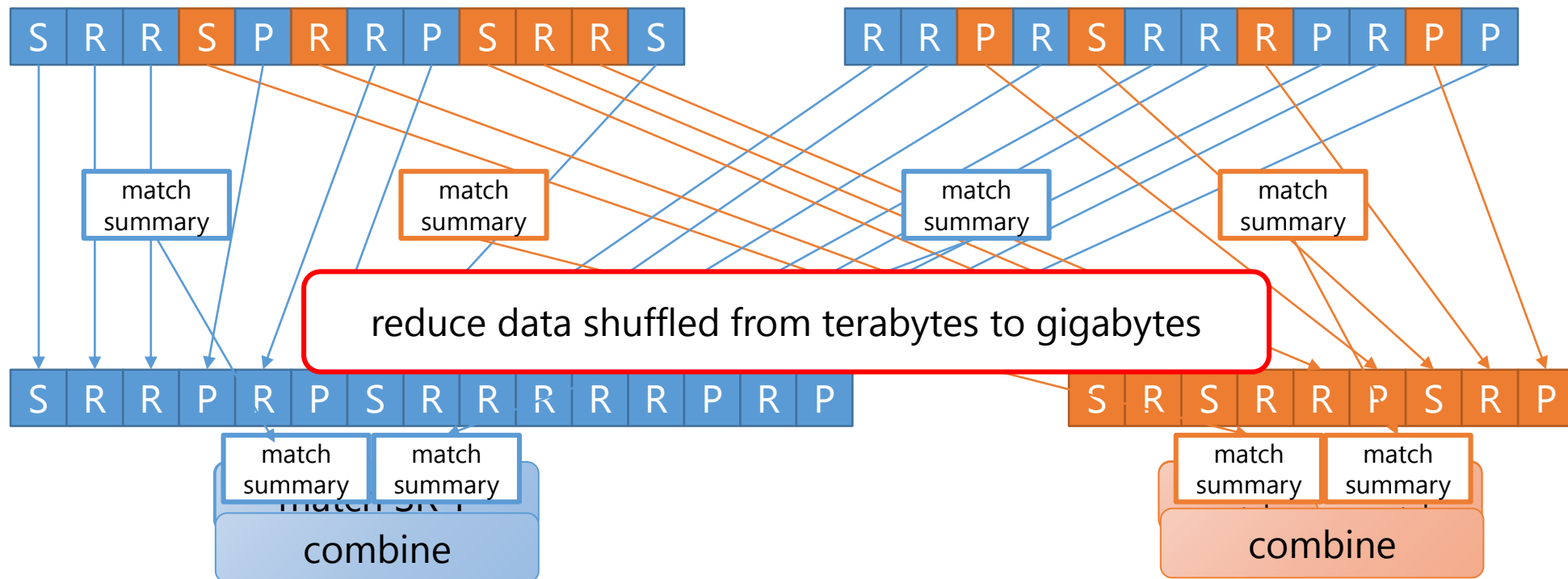


# Count the number of reviews read per user





# Count influential reviews ( $SR^+P$ ) per user



# SymPLE [SOSP '15]

a language for specifying nonrelational parts of data-processing queries  
a subset of C++

automatically parallelize sequential code

expose additional parallelism to query optimizer

up to 2 orders of magnitude efficiency improvement

# Count influential reviews

```
bool search_done = false;
int num_reviews = 0;
int sum = 0;

for each record in input
    switch record.type:
        case SEARCH:    if (!search_done) { num_reviews = 0;
                                search_done = true; }

        case REVIEW:    num_reviews++;

        case PURCHASE:  if (search_done) { search_done = false;
                                sum += num_reviews; }
```

# Count influential reviews

```
SymBool search_done = false;  
SymInt num_reviews = 0;  
SymInt sum = 0;
```

user uses symbolic data types  
for loop carried state

```
for each record in input  
  switch record.type:  
    case SEARCH:    if (!search_done) { num_reviews = 0;  
                                     search_done = true; }  
    case REVIEW:    num_reviews++;  
    case PURCHASE:  if (search_done) { search_done = false;  
                                     sum += num_reviews; }
```

overloaded operators encode  
efficient symbolic decision  
procedures for generating  
symbolic summaries

# Computing max in parallel

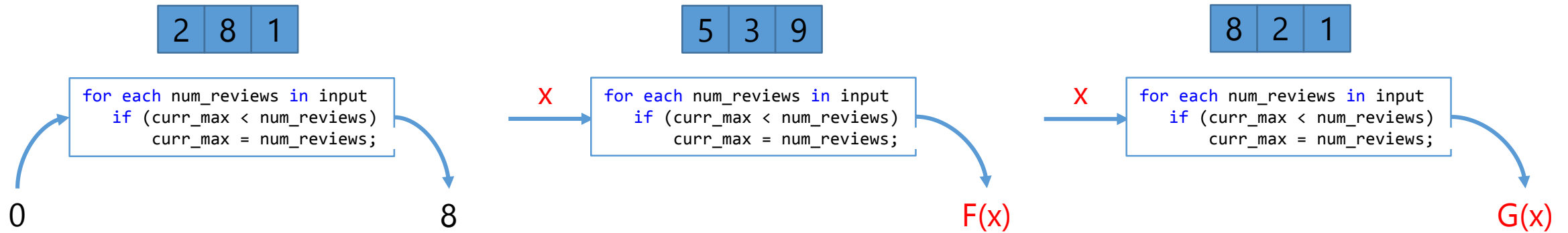
```
SymInt curr_max = 0;  
  
for each num_reviews in input  
  if (curr_max < num_reviews)  
    curr_max = num_reviews;
```

max is, of course, associative

but this is not apparent from code

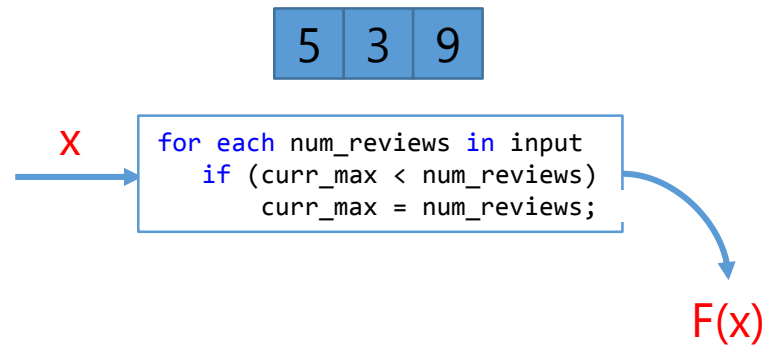
SymPLE can parallelize this code

# Parallelize by breaking dependences



output = G(F(8))

# Parallelize by breaking dependences



```

SymInt max = x;
for each num_reviews in (5,3,9)
  if (max < num_reviews)
    max = num_reviews;

```

no branching  
when state  
becomes  
concrete

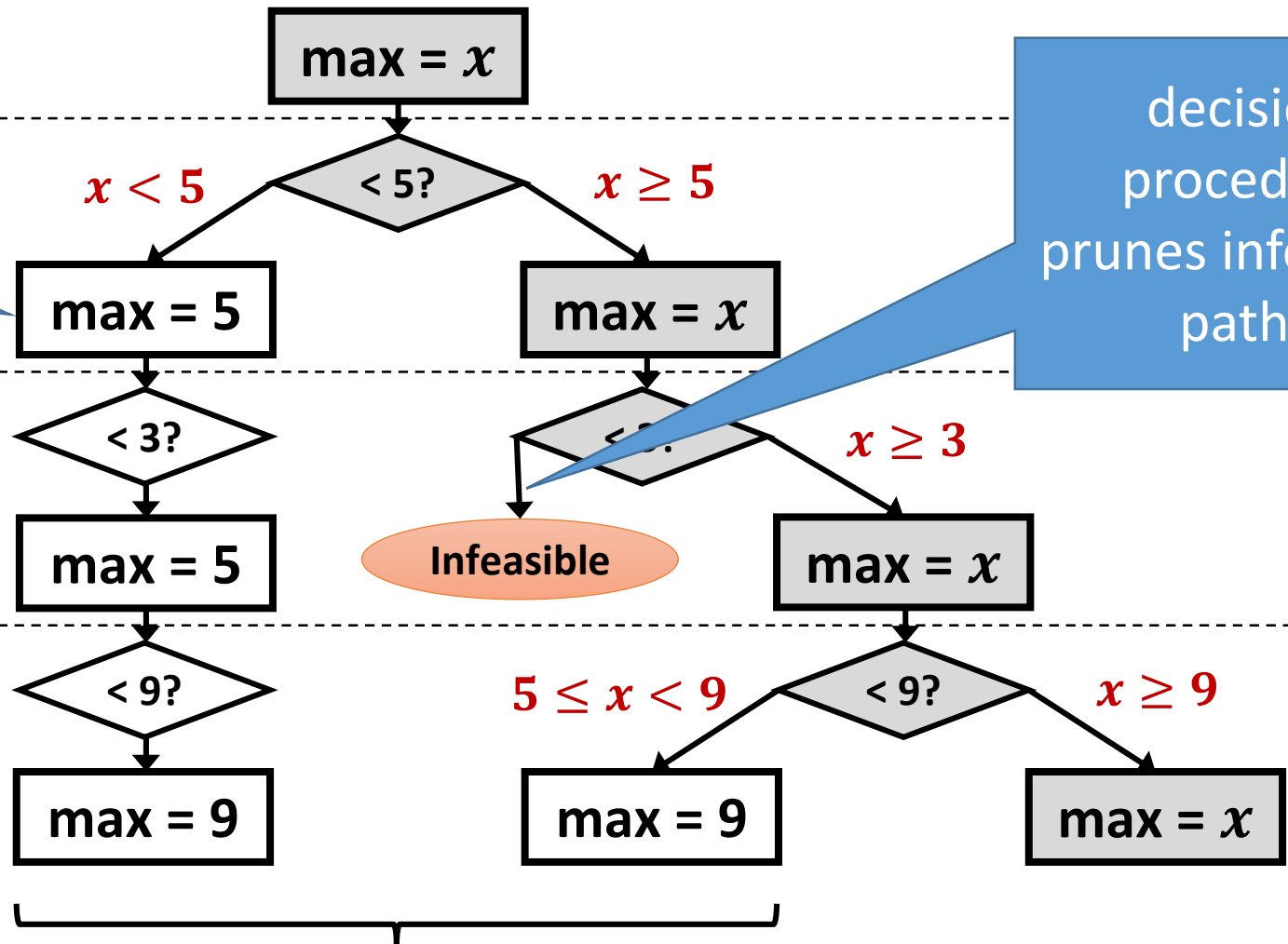
decision  
procedure  
prunes infeasible  
paths

```

if (max < 3)
  max = 3;

```

equivalent  
paths can be  
merged

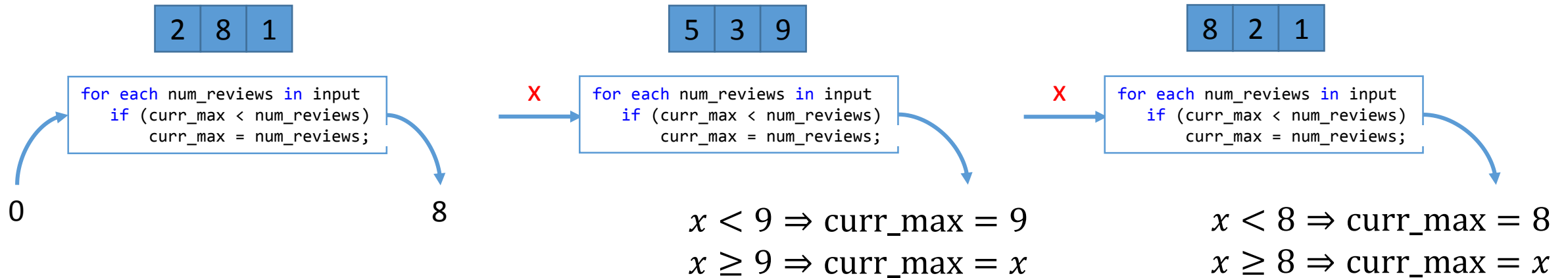


$x < 9 \Rightarrow \text{max} = 9$

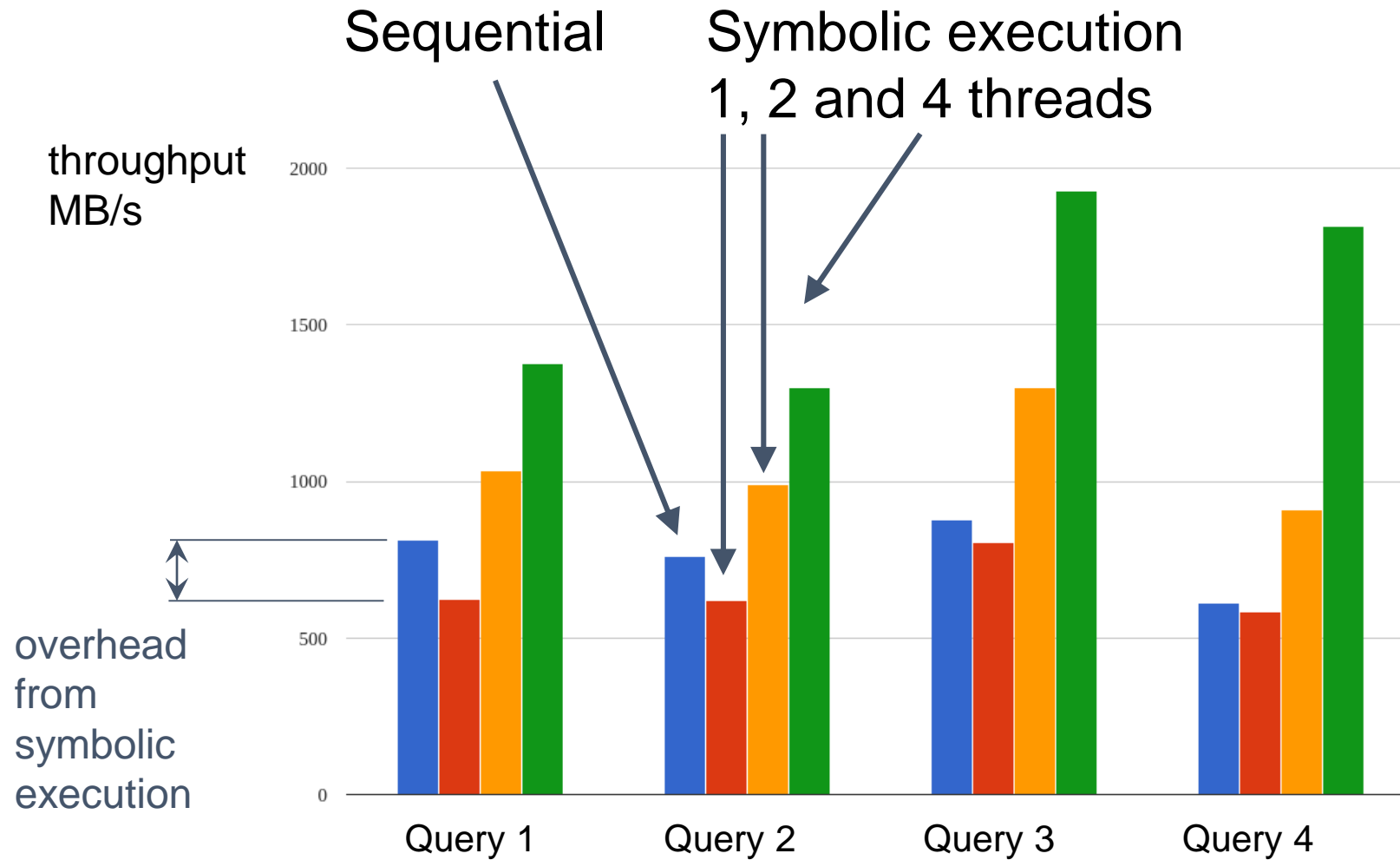
$x \geq 9 \Rightarrow \text{max} = x$



# Parallelize by breaking dependences



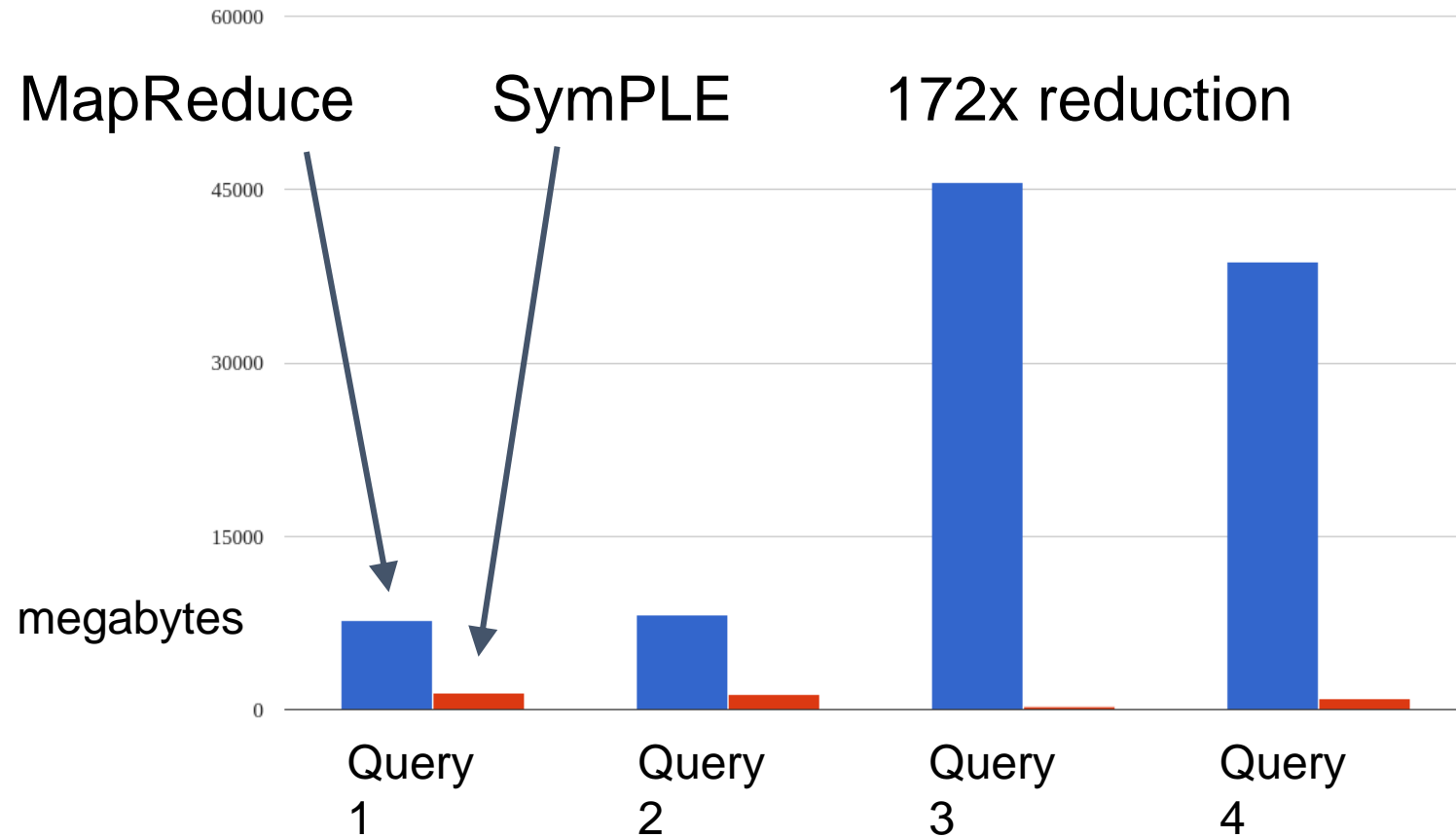
# Single machine throughput



# Reduction in data movement



data shuffled from mappers to reducers



# Challenge

can we develop new abstractions for future data-processing needs?

- move beyond embarrassingly parallel
- automatically parallelizable

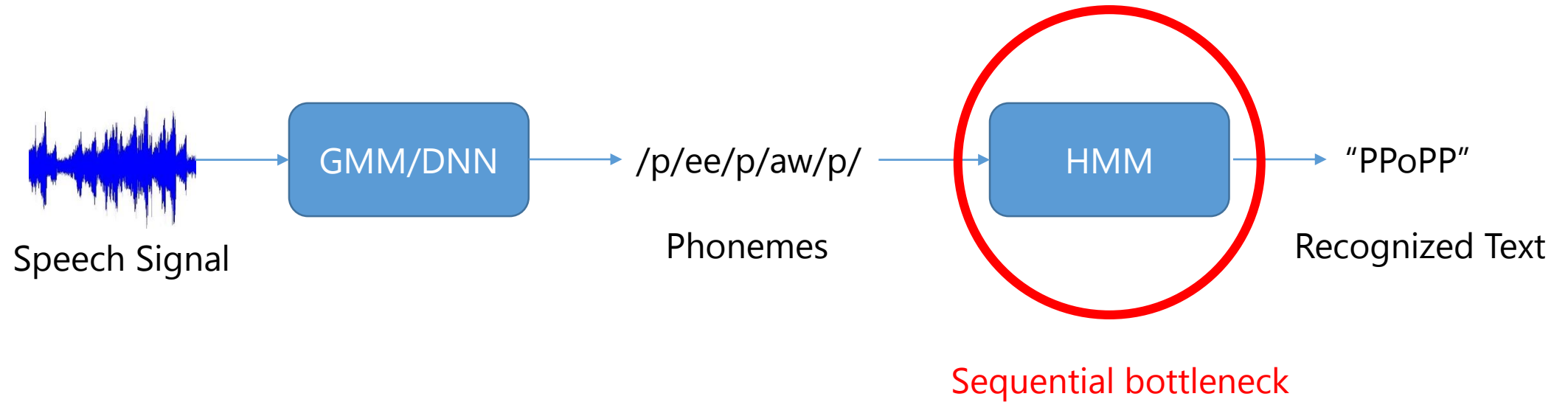
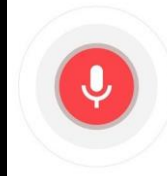
perform whole query optimizations

- unify relational and non-relational parts
- extract filters, project unused parts of data, ...

# Manual Parallelization Across Dependences

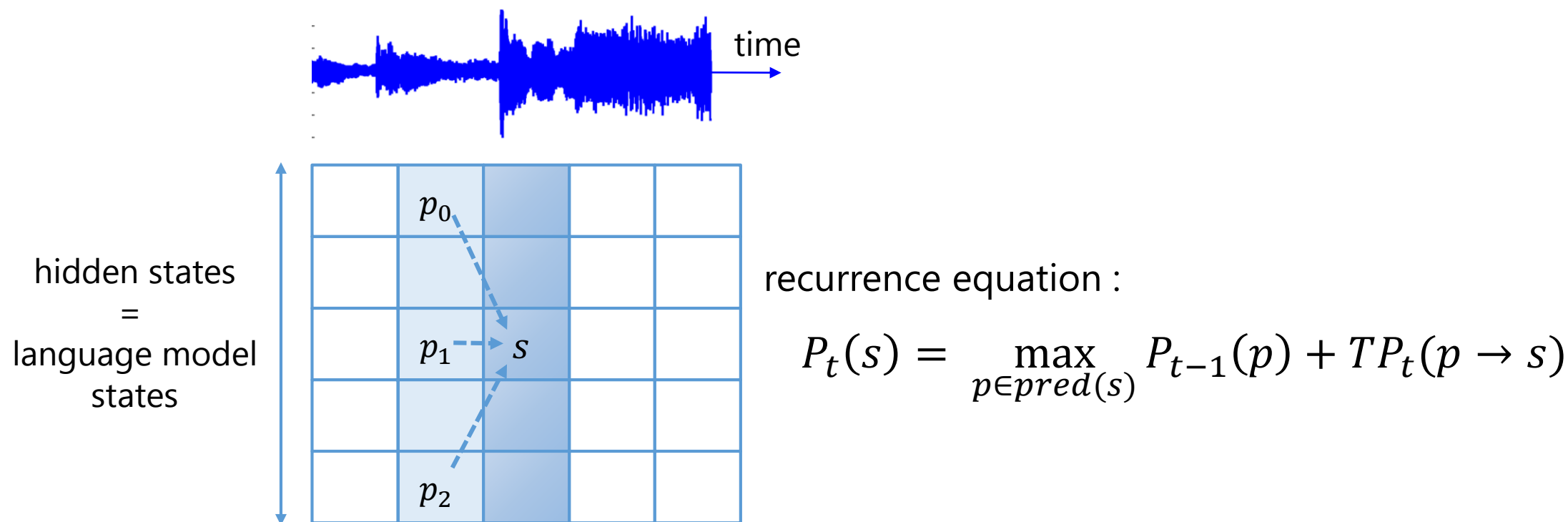
## Dynamic Programming

# Speech decoders



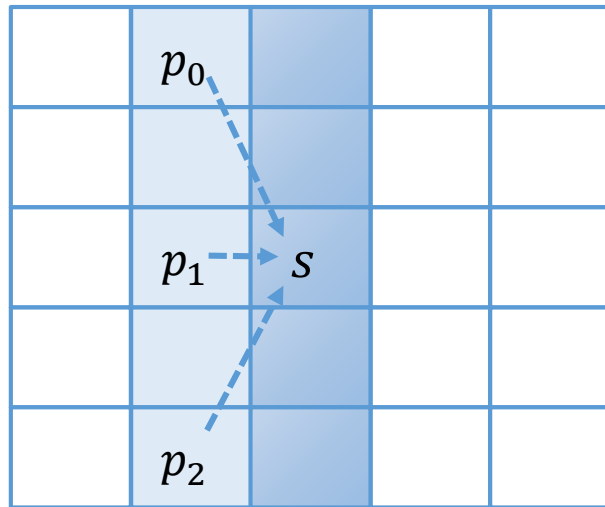
# Viterbi algorithm for Hidden Markov Models (HMM)

finds the most likely sequence of hidden states that explain an observation



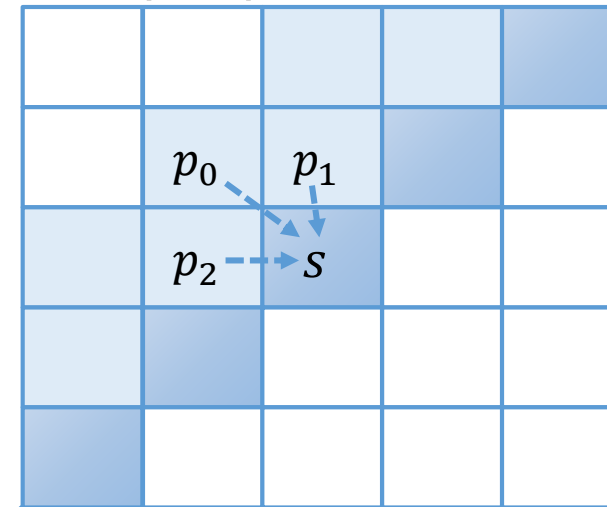
# Dynamic programming computes a sequence of stages

Viterbi



stage = column

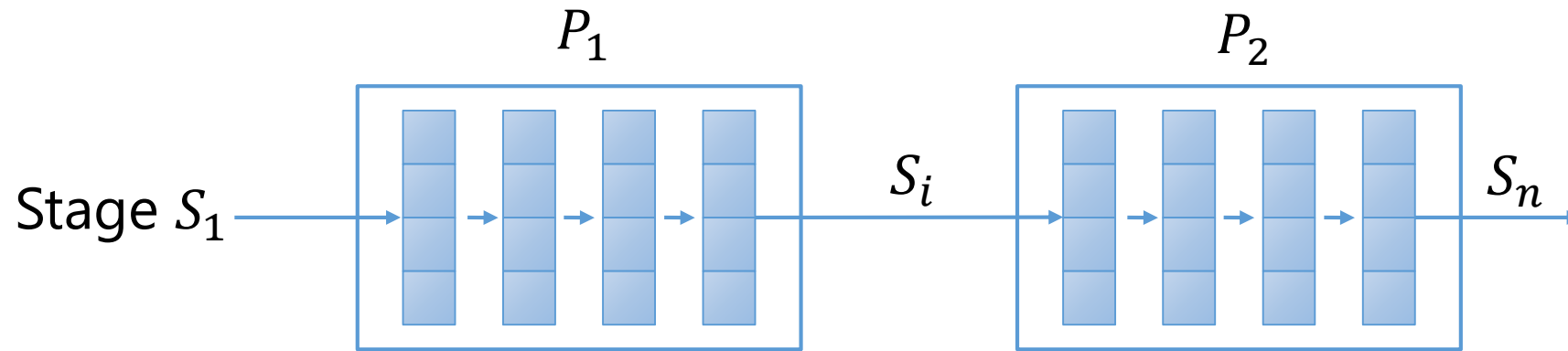
LCS (diff)



stage = anti-diagonal



# Our focus: parallelization across stages



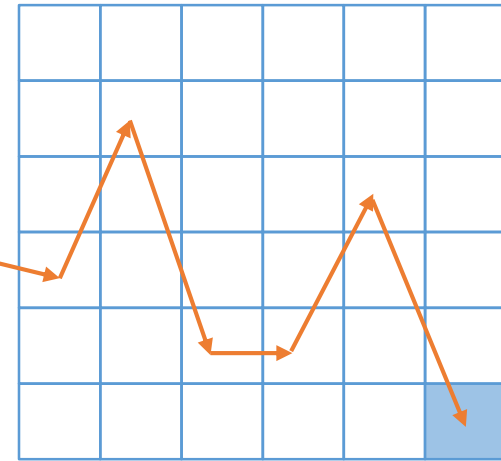
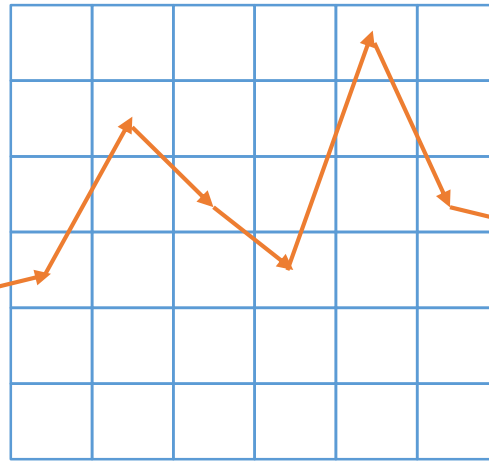
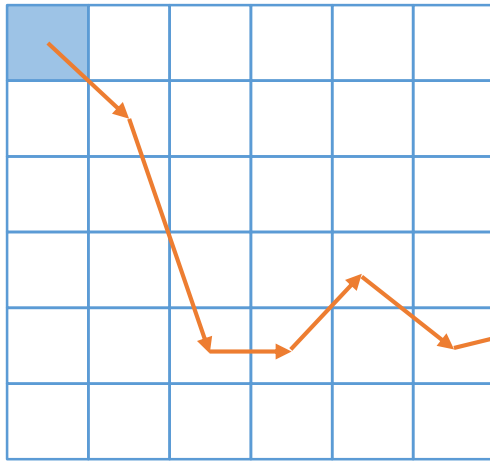
$$S_t[i] = \max_j (S_{t-1}[j] + c_{t,i,j})$$

$$\overrightarrow{S_t} = A \odot \overrightarrow{S_{t-1}}$$

where  $\odot$  is matrix multiplication in tropical semiring

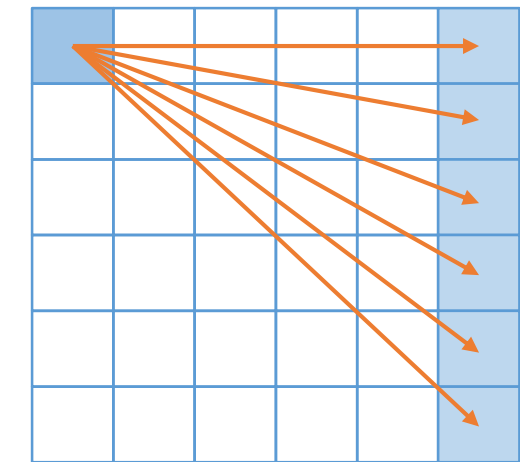
# Solution in terms of finding shortest-paths

source

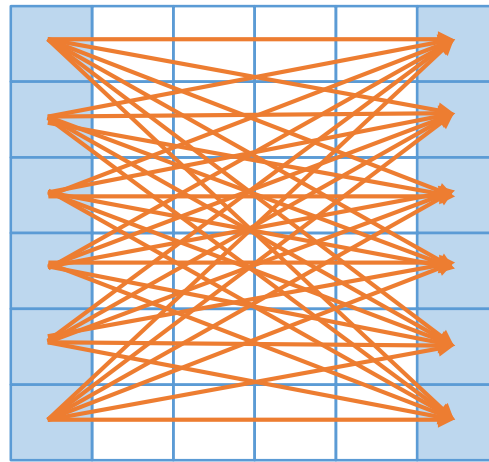


dest

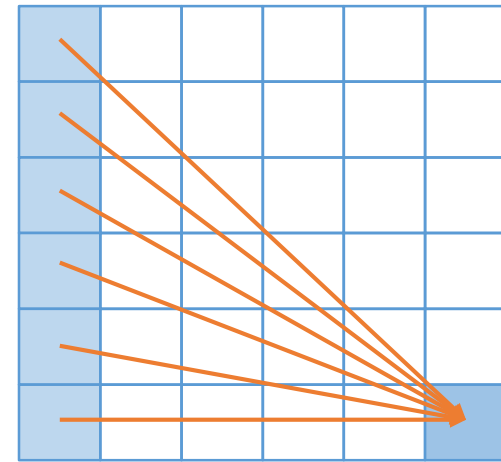
# Solution in terms of finding shortest-paths



src → all dest



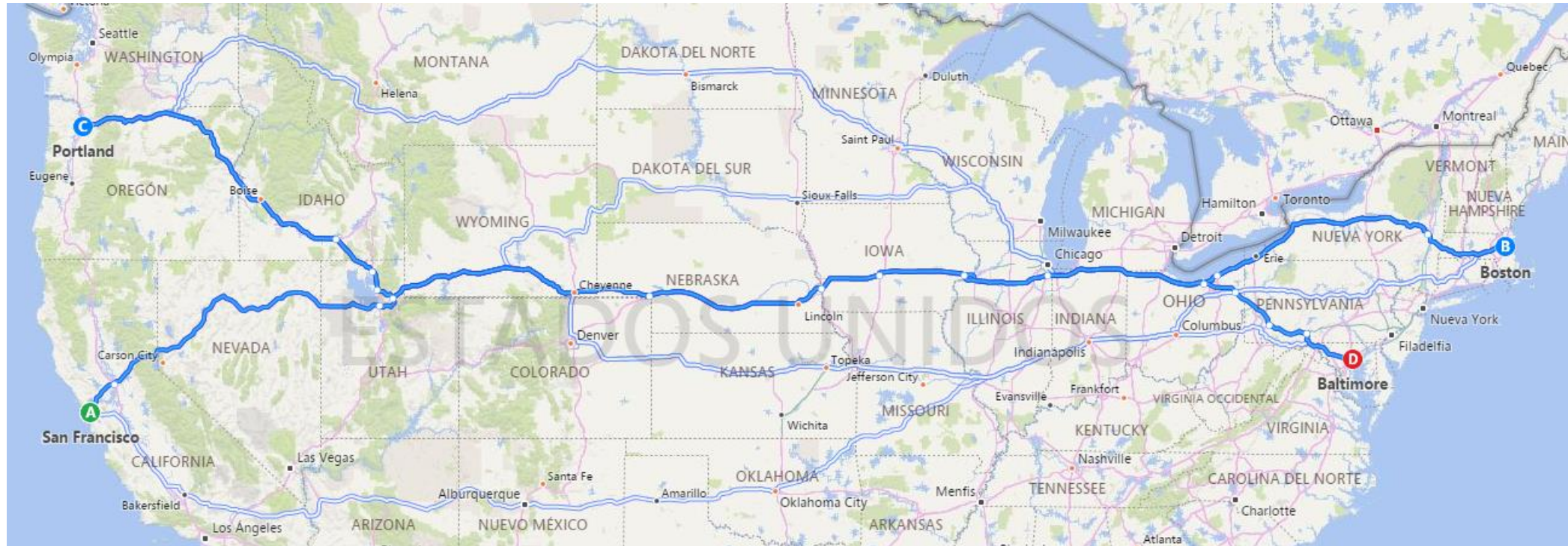
all src → all dest



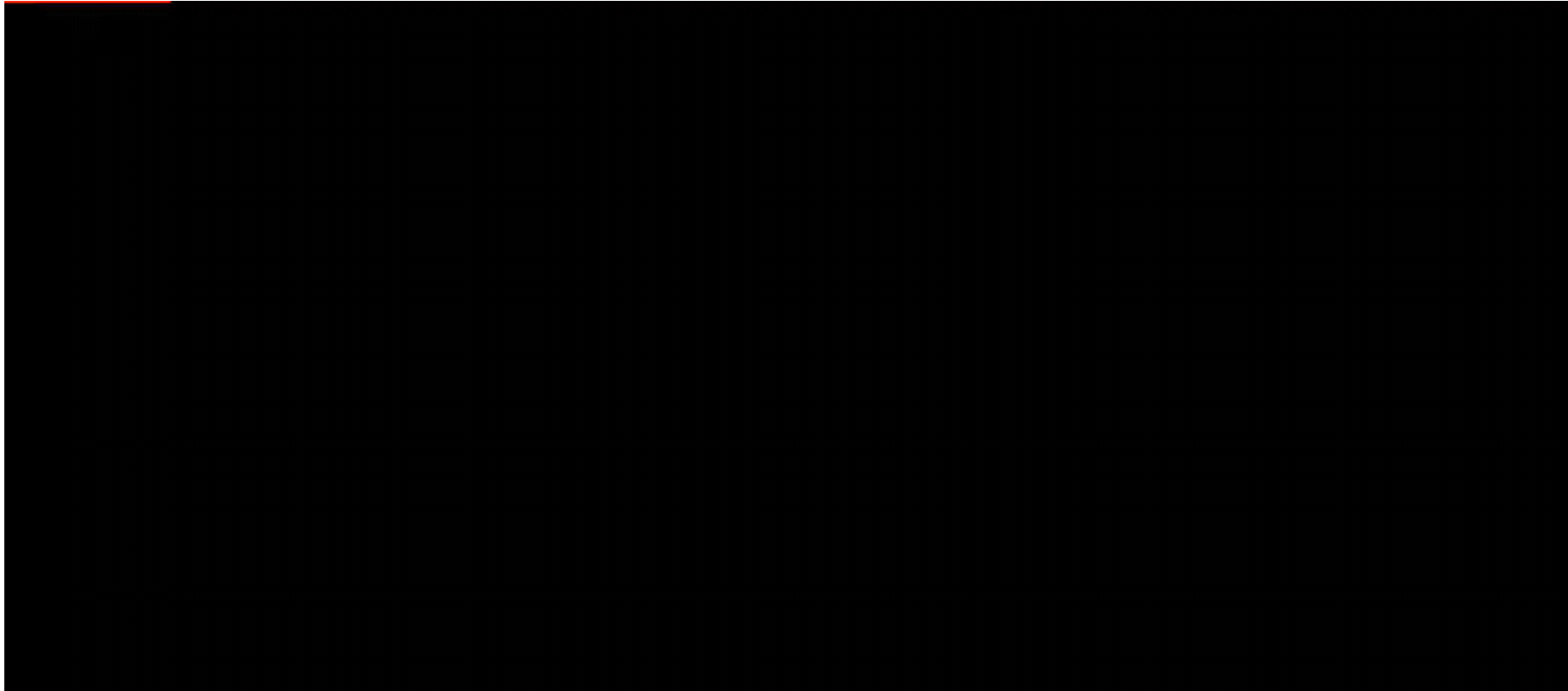
all src → dest

parallelization cost = size of stages

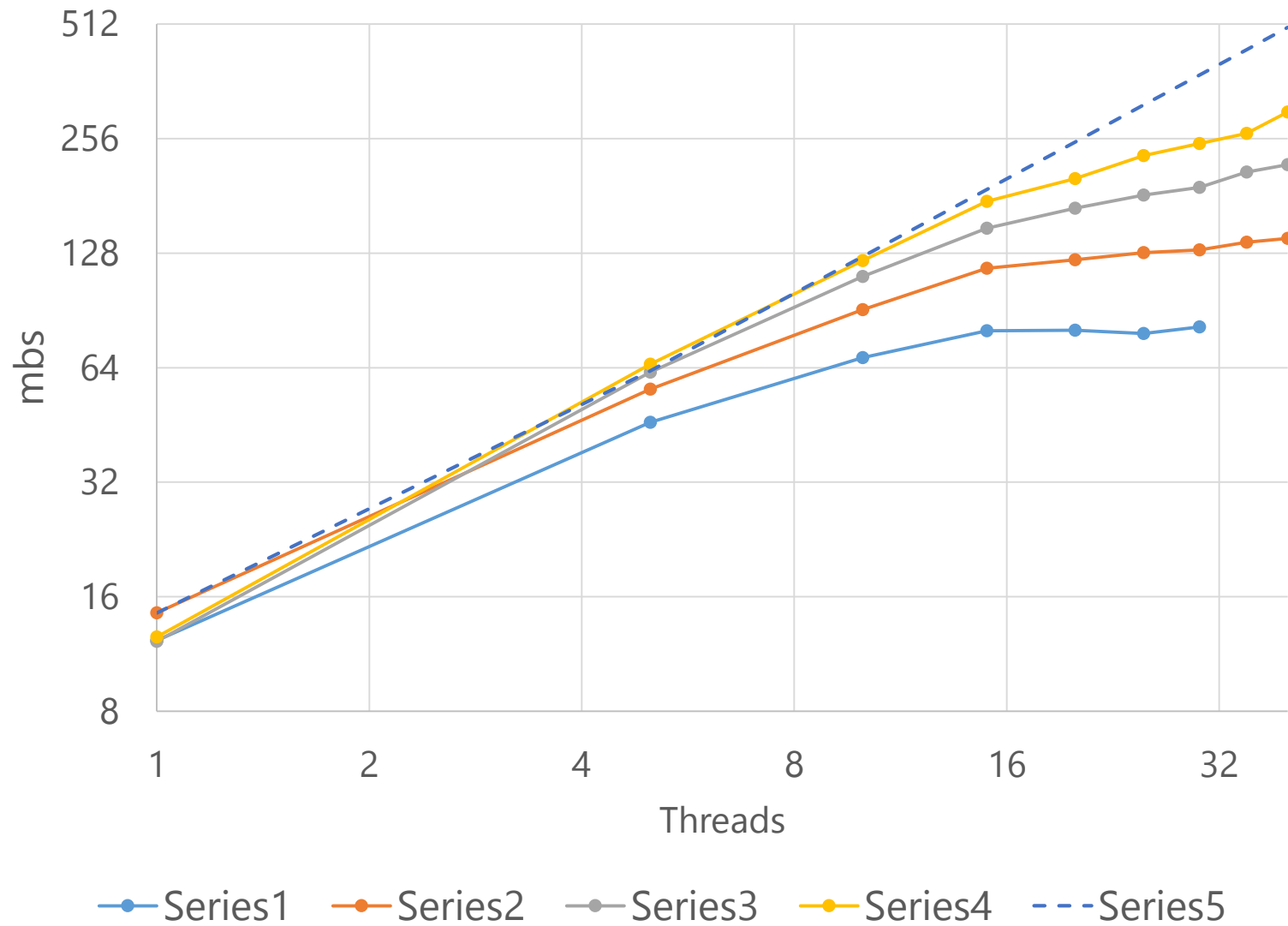
# Shortest paths converge to optimal routes



# Convergence in LCS



# Speed of Viterbi Decoder on CDMA



# Summary

“inherently sequential”  $\Rightarrow$  “embarrassingly parallel”

